

On global algebraic completeness of the Gödel-Löb provability logic

Daniyar Shamkanov
daniyar.shamkanov@gmail.com

Steklov Mathematical Institute of Russian Academy of Sciences
National Research University Higher School of Economics
Moscow

TACL, Nice
June 17-21, 2019

Gödel-Löb provability logic GL

Axiom schemes

- ▶ Boolean tautologies
- ▶ $\Box(A \rightarrow B) \rightarrow (\Box A \rightarrow \Box B)$
- ▶ $\Box(\Box A \rightarrow A) \rightarrow \Box A$ (Löb's axiom)

Inference rules

$$\text{mp} \frac{A \quad A \rightarrow B}{B}, \quad \text{nec} \frac{A}{\Box A}.$$

The logic GL is sound and complete w.r.t. the arithmetical semantics, where $\Box A$ corresponds to "A is provable in Peano arithmetic".

Non-well-founded derivations

Definition

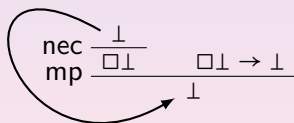
An ∞ -*derivation* in GL is a (possibly infinite) tree whose nodes are marked by modal formulas and that is constructed according to the rules (mp) and (nec). In addition, any infinite branch in an ∞ -derivation must contain infinitely many applications of the rule (nec).

Non-well-founded derivations

Definition

An ∞ -*derivation* in GL is a (possibly infinite) tree whose nodes are marked by modal formulas and that is constructed according to the rules (mp) and (nec). In addition, any infinite branch in an ∞ -derivation must contain infinitely many applications of the rule (nec).

Example

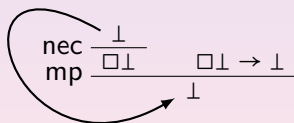


Non-well-founded derivations

Definition

An ∞ -*derivation* in GL is a (possibly infinite) tree whose nodes are marked by modal formulas and that is constructed according to the rules (mp) and (nec). In addition, any infinite branch in an ∞ -derivation must contain infinitely many applications of the rule (nec).

Example



Theorem

GL = K4 + cyclic proofs = K4 + ∞ -proofs

Non-well-founded derivations

Definition

An ∞ -*derivation* in GL is a (possibly infinite) tree whose nodes are marked by modal formulas and that is constructed according to the rules (mp) and (nec). In addition, any infinite branch in an ∞ -derivation must contain infinitely many applications of the rule (nec).

Example

$$\begin{array}{c} \vdots \\ \text{mp} \frac{\Box p_3 \quad \Box p_3 \rightarrow p_2}{p_2} \\ \text{nec} \frac{p_2}{\Box p_2} \\ \text{mp} \frac{\Box p_2 \quad \Box p_2 \rightarrow p_1}{p_1} \\ \text{nec} \frac{p_1}{\Box p_1} \\ \text{mp} \frac{\Box p_1 \quad \Box p_1 \rightarrow p_0}{p_0} \end{array}$$

Non-well-founded derivations

Definition

An *assumption leaf* of an ∞ -derivation is a leaf that is not marked by an axiom of GL. An assumption leaf is *boxed* if there is an application of the rule (nec) on the path from this leaf to the root of the tree.

Example

$$\begin{array}{c} \vdots \\ \text{mp} \frac{\boxed{p_3} \quad \boxed{p_3 \rightarrow p_2}}{p_2} \\ \text{nec} \frac{p_2}{\boxed{p_2}} \\ \text{mp} \frac{\boxed{p_2} \quad \boxed{p_2 \rightarrow p_1}}{p_1} \\ \text{nec} \frac{p_1}{\boxed{p_1}} \\ \text{mp} \frac{\boxed{p_1} \quad \boxed{p_1 \rightarrow p_0}}{p_0} \end{array}$$

Non-well-founded derivations

Definition

An *assumption leaf* of an ∞ -derivation is a leaf that is not marked by an axiom of GL. An assumption leaf is *boxed* if there is an application of the rule (nec) on the path from this leaf to the root of the tree.

Example

$$\begin{array}{c} \vdots \\ \text{mp} \frac{\boxed{p_3} \quad \boxed{p_3 \rightarrow p_2}}{p_2} \\ \text{nec} \frac{p_2}{\boxed{p_2}} \\ \text{mp} \frac{\boxed{p_2} \quad \boxed{p_2 \rightarrow p_1}}{p_1} \\ \text{nec} \frac{p_1}{\boxed{p_1}} \\ \text{mp} \frac{\boxed{p_1} \quad \boxed{p_1 \rightarrow p_0}}{p_0} \end{array}$$

Local and global derivability relations

Definition

We set $\Gamma \vdash_l A$ if there is a finite derivation of A from assumptions Γ and theorems of GL constructed according to the rule (mp).

Local and global derivability relations

Definition

We set $\Gamma \vdash_l A$ if there is a finite derivation of A from assumptions Γ and theorems of GL constructed according to the rule (mp).

Definition

We set $\Gamma \vdash_g A$ if there is an ∞ -derivation with the root marked by A in which all assumption leafs are marked by some elements of Γ .

Local and global derivability relations

Definition

We set $\Gamma \vdash_l A$ if there is a finite derivation of A from assumptions Γ and theorems of GL constructed according to the rule (mp).

Definition

We set $\Gamma \vdash_g A$ if there is an ∞ -derivation with the root marked by A in which all assumption leafs are marked by some elements of Γ .

Definition

We also set $\Sigma; \Gamma \vdash A$ if there is an ∞ -derivation with the root marked by A in which all boxed assumption leafs are marked by some elements of Σ and all non-boxed assumption leafs are marked by some elements of Γ .

Local and global derivability relations

Definition

We set $\Gamma \vdash_l A$ if there is a finite derivation of A from assumptions Γ and theorems of GL constructed according to the rule (mp).

Definition

We set $\Gamma \vdash_g A$ if there is an ∞ -derivation with the root marked by A in which all assumption leafs are marked by some elements of Γ .

Definition

We also set $\Sigma; \Gamma \vdash A$ if there is an ∞ -derivation with the root marked by A in which all boxed assumption leafs are marked by some elements of Σ and all non-boxed assumption leafs are marked by some elements of Γ .

The relation \vdash is a generalization of \vdash_l and \vdash_g :

$$\Gamma \vdash_l A \iff \emptyset; \Gamma \vdash A \quad \text{and} \quad \Gamma \vdash_g A \iff \Gamma; \Gamma \vdash A.$$

Global neighbourhood completeness of GL

Definition

We set $\Gamma \vDash_g A$ if for any neighbourhood GL-model \mathcal{M}

$$(\forall B \in \Gamma \mathcal{M} \vDash B) \Rightarrow \mathcal{M} \vDash A.$$

Global neighbourhood completeness of GL

Definition

We set $\Gamma \vDash_g A$ if for any neighbourhood GL-model \mathcal{M}

$$(\forall B \in \Gamma \mathcal{M} \vDash B) \Rightarrow \mathcal{M} \vDash A.$$

Theorem

$$\Gamma \vdash_g A \iff \Gamma \vDash_g A.$$

Algebraic semantics

Magari algebras

Definition

A *Magari algebra* $\mathcal{A} = (X, \wedge, \vee, \rightarrow, 0, 1, \square)$ is a Boolean algebra $(X, \wedge, \vee, \rightarrow, 0, 1)$ together with a unary map $\square: X \rightarrow X$ satisfying the identities:

$$\square 1 = 1, \quad \square(x \wedge y) = \square x \wedge \square y, \quad \square(\square x \rightarrow x) = \square x.$$

Magari algebras

Definition

A *Magari algebra* $\mathcal{A} = (X, \wedge, \vee, \rightarrow, 0, 1, \Box)$ is a Boolean algebra $(X, \wedge, \vee, \rightarrow, 0, 1)$ together with a unary map $\Box: X \rightarrow X$ satisfying the identities:

$$\Box 1 = 1, \quad \Box(x \wedge y) = \Box x \wedge \Box y, \quad \Box(\Box x \rightarrow x) = \Box x.$$

Definition

A Magari algebra \mathcal{A} is *\Box -founded* if, for every sequence of elements $(b_i)_{i \in \mathbb{N}}$ such that $\Box b_{i+1} \leq b_i$, we have $b_0 = 1$.

Definition

A Magari algebra is σ -complete if its underlying Boolean algebra is σ -complete that is every its countable subset S has the least upper bound $\bigvee S$ (equivalently, the greatest lower bound $\bigwedge S$).

Proposition

Any σ -complete Magari algebra is \square -founded.

Definition

A Magari algebra is σ -complete if its underlying Boolean algebra is σ -complete that is every its countable subset S has the least upper bound $\bigvee S$ (equivalently, the greatest lower bound $\bigwedge S$).

Proposition

Any σ -complete Magari algebra is \square -founded.

Proof

Assume we have a σ -complete Magari algebra \mathcal{A} and a sequence of its elements $(a_i)_{i \in \mathbb{N}}$ such that $\square a_{i+1} \leq a_i$. We shall prove that $a_0 = 1$. Put $b = \bigwedge_{i \in \mathbb{N}} a_i$. For any $n \in \mathbb{N}$, we have $b \leq a_{n+1}$ and $\square b \leq \square a_{n+1} \leq a_n$. Hence,

$$\square b \leq b, \quad \square b \rightarrow b = 1, \quad \square b = \square(\square b \rightarrow b) = 1, \quad b = 1.$$

We obtain that $a_0 = 1$.

Algebraic consequence relations

A *valuation on \mathcal{A}* is a function $\theta: Fm \rightarrow X$ such that $\theta(\perp) = 0$, $\theta(A \rightarrow B) = \theta(A) \rightarrow \theta(B)$ and $\theta(\Box A) = \Box\theta(A)$. For a subset S of a Magari algebra \mathcal{A} , the filter of (the boolean part of) \mathcal{A} generated by S is denoted by $\langle S \rangle$.

Definition

Given a set of modal formulas Γ and a formula A , we set $\Gamma \models_I A$ if for any \Box -founded Magari algebra \mathcal{A} and any valuation θ on \mathcal{A}

$$\theta(A) \in \langle \{ \theta(B) \mid B \in \Gamma \} \rangle.$$

Definition

We also set $\Gamma \models_g A$ if for any \Box -founded Magari algebra \mathcal{A} and any valuation θ on \mathcal{A}

$$(\forall B \in \Gamma \theta(B) = 1) \Rightarrow \theta(A) = 1.$$

Algebraic consequence relations

Definition

We set $\Sigma; \Gamma \models A$ if for any \square -founded Magari algebra \mathcal{A} and any valuation θ on \mathcal{A}

$$(\forall C \in \Sigma \ \square \theta(C) = 1) \Rightarrow \theta(A) \in \langle \{\theta(B) \mid B \in \Gamma\} \rangle.$$

Algebraic consequence relations

Definition

We set $\Sigma; \Gamma \Vdash A$ if for any \square -founded Magari algebra \mathcal{A} and any valuation θ on \mathcal{A}

$$(\forall C \in \Sigma \ \square \theta(C) = 1) \Rightarrow \theta(A) \in \langle \{\theta(B) \mid B \in \Gamma\} \rangle.$$

The relation \Vdash is a generalization of \Vdash_I and \Vdash_g :

$$\Gamma \Vdash_I A \iff \emptyset; \Gamma \Vdash A \quad \text{and} \quad \Gamma \Vdash_g A \iff \Gamma; \Gamma \Vdash A.$$

Algebraic consequence relations

Definition

We set $\Sigma; \Gamma \models A$ if for any \square -founded Magari algebra \mathcal{A} and any valuation θ on \mathcal{A}

$$(\forall C \in \Sigma \ \square \theta(C) = 1) \Rightarrow \theta(A) \in \langle \{\theta(B) \mid B \in \Gamma\} \rangle.$$

The relation \models is a generalization of \models_I and \models_g :

$$\Gamma \models_I A \iff \emptyset; \Gamma \models A \quad \text{and} \quad \Gamma \models_g A \iff \Gamma; \Gamma \models A.$$

Theorem

$$\Sigma; \Gamma \models A \iff \Sigma; \Gamma \vdash A.$$

Global neighbourhood completeness of GL is essentially global completeness over atomic complete Magari algebras.

Global neighbourhood completeness of GL is essentially global completeness over atomic complete Magari algebras.

Plan of the proof

- ▶ Correctness
- ▶ Compactness
- ▶ Completeness via a sequent calculus

Completeness via a sequent calculus

Definition

A *sequent* is an expression of the form $\Sigma; \Gamma \Rightarrow \Delta$, where Γ and Δ are finite multisets of formulas, and Σ is an arbitrary set of formulas.

Definition

A sequent $\Sigma; \Gamma \Rightarrow \Delta$ is called *valid* if $\Sigma; \{\wedge \Gamma\} \models \vee \Delta$.

Non-well-founded sequent calculus

For a finite multiset of formulas $\Gamma = A_1, \dots, A_n$, we set $\Box\Gamma := \Box A_1, \dots, \Box A_n$.

Initial sequents and inference rules of the sequent calculus S have the following form:

$$\Sigma; \Gamma, p \Rightarrow p, \Delta, \quad \Sigma; \Gamma, \perp \Rightarrow \Delta,$$

$$\rightarrow_L \frac{\Sigma; \Gamma, B \Rightarrow \Delta \quad \Sigma; \Gamma \Rightarrow A, \Delta}{\Sigma; \Gamma, A \rightarrow B \Rightarrow \Delta}, \quad \rightarrow_R \frac{\Sigma; \Gamma, A \Rightarrow B, \Delta}{\Sigma; \Gamma \Rightarrow A \rightarrow B, \Delta},$$

$$\Box \frac{\Sigma; \Sigma_0, \Gamma, \Box\Gamma \Rightarrow A}{\Sigma; \Pi, \Box\Gamma \Rightarrow \Box A, \Delta} \quad (\Sigma_0 \text{ is a finite subset of } \Sigma).$$

Lemma

If a sequent $\Sigma; \Gamma, A \rightarrow B \Rightarrow \Delta$ is valid, then sequents $\Sigma; \Gamma, B \Rightarrow \Delta$ and $\Sigma; \Gamma \Rightarrow A, \Delta$ are valid. If $\Sigma; \Gamma \Rightarrow A \rightarrow B, \Delta$ is valid, then $\Sigma; \Gamma, A \Rightarrow B, \Delta$ is also valid.

Definition

A sequent $\Sigma; \Gamma \Rightarrow \Delta$ is called *saturated* if Γ and Δ do not contain formulas of the form $A \rightarrow B$.

Lemma

If $\Sigma; \Pi, \Box\Gamma \Rightarrow \Delta$ is a valid non-initial saturated sequent, where Π consists only of propositional variables, then there are a finite subset Σ_0 of Σ and a formula $\Box A$ from Δ such that $\Sigma; \Sigma_0, \Gamma, \Box\Gamma \Rightarrow A$ is a valid sequent.

Initial sequents and inference rules of the sequent calculus S:

$$\Sigma; \Gamma, p \Rightarrow p, \Delta, \quad \Sigma; \Gamma, \perp \Rightarrow \Delta,$$

$$\rightarrow_L \frac{\Sigma; \Gamma, B \Rightarrow \Delta \quad \Sigma; \Gamma \Rightarrow A, \Delta}{\Sigma; \Gamma, A \rightarrow B \Rightarrow \Delta}, \quad \rightarrow_R \frac{\Sigma; \Gamma, A \Rightarrow B, \Delta}{\Sigma; \Gamma \Rightarrow A \rightarrow B, \Delta},$$

$$\Box \frac{\Sigma; \Sigma_0, \Gamma, \Box \Gamma \Rightarrow A}{\Sigma; \Pi, \Box \Gamma \Rightarrow \Box A, \Delta} \quad (\Sigma_0 \text{ is a finite subset of } \Sigma).$$

Initial sequents and inference rules of the sequent calculus S:

$$\Sigma; \Gamma, p \Rightarrow p, \Delta, \quad \Sigma; \Gamma, \perp \Rightarrow \Delta,$$

$$\rightarrow_L \frac{\Sigma; \Gamma, B \Rightarrow \Delta \quad \Sigma; \Gamma \Rightarrow A, \Delta}{\Sigma; \Gamma, A \rightarrow B \Rightarrow \Delta}, \quad \rightarrow_R \frac{\Sigma; \Gamma, A \Rightarrow B, \Delta}{\Sigma; \Gamma \Rightarrow A \rightarrow B, \Delta},$$

$$\Box \frac{\Sigma; \Sigma_0, \Gamma, \Box \Gamma \Rightarrow A}{\Sigma; \Pi, \Box \Gamma \Rightarrow \Box A, \Delta} \quad (\Sigma_0 \text{ is a finite subset of } \Sigma).$$

Theorem

Any valid sequent is provable in S.

Initial sequents and inference rules of the sequent calculus S:

$$\Sigma; \Gamma, p \Rightarrow p, \Delta, \quad \Sigma; \Gamma, \perp \Rightarrow \Delta,$$

$$\rightarrow_L \frac{\Sigma; \Gamma, B \Rightarrow \Delta \quad \Sigma; \Gamma \Rightarrow A, \Delta}{\Sigma; \Gamma, A \rightarrow B \Rightarrow \Delta}, \quad \rightarrow_R \frac{\Sigma; \Gamma, A \Rightarrow B, \Delta}{\Sigma; \Gamma \Rightarrow A \rightarrow B, \Delta},$$

$$\Box \frac{\Sigma; \Sigma_0, \Gamma, \Box \Gamma \Rightarrow A}{\Sigma; \Pi, \Box \Gamma \Rightarrow \Box A, \Delta} \quad (\Sigma_0 \text{ is a finite subset of } \Sigma).$$

Theorem

Any valid sequent is provable in S.

Corollary (Algebraic completeness)

$$\Sigma; \Gamma \models A \Rightarrow \Sigma; \Gamma \vdash A.$$

Global neighbourhood completeness of GL is essentially global completeness over atomic complete Magari algebras.

Plan of the proof

- ▶ Correctness
- ▶ Compactness
- ▶ Completeness via a sequent calculus

Global neighbourhood completeness of GL is essentially global completeness over atomic complete Magari algebras.

Plan of the proof

- ▶ Correctness
- ▶ Compactness
- ▶ Completeness via a sequent calculus

Theorem

A Magari algebra is \square -founded if and only if it is embeddable into an atomic complete Magari algebra.

Global completeness of GL over σ -complete Magari algebras can be obtained along the same lines.

Plan of the proof

- ▶ Correctness
- ▶ Compactness
- ▶ Completeness via a sequent calculus

Global completeness of GL over σ -complete Magari algebras can be obtained along the same lines.

Plan of the proof

- ▶ Correctness
- ▶ Compactness
- ▶ Completeness via a sequent calculus

Proposition

A Magari algebra is \square -founded if and only if it is embeddable into a σ -complete Magari algebra.

Consequence relation over σ -complete Magari algebras

Definition

We also set $\Gamma \vDash_g A$ if for any σ -complete Magari algebra \mathcal{A} and any valuation θ on \mathcal{A}

$$(\forall B \in \Gamma \theta(B) = 1) \Rightarrow \theta(A) = 1.$$

Definition

We set $\Sigma; \Gamma \vDash A$ if for any σ -complete Magari algebra \mathcal{A} and any valuation θ in \mathcal{A}

$$\square \bigwedge \{ \theta(C) \mid C \in \Sigma \} = 1 \Rightarrow \theta(A) \in \langle \{ \theta(B) \mid B \in \Gamma \} \rangle.$$

Consequence relation over σ -complete Magari algebras

Definition

We also set $\Gamma \vDash_g A$ if for any σ -complete Magari algebra \mathcal{A} and any valuation θ on \mathcal{A}

$$(\forall B \in \Gamma \theta(B) = 1) \Rightarrow \theta(A) = 1.$$

Definition

We set $\Sigma; \Gamma \vDash A$ if for any σ -complete Magari algebra \mathcal{A} and any valuation θ in \mathcal{A}

$$\square \bigwedge \{\theta(C) \mid C \in \Sigma\} = 1 \Rightarrow \theta(A) \in \langle \{\theta(B) \mid B \in \Gamma\} \rangle.$$

Lemma

$\Sigma; \Gamma \Vdash A \Rightarrow \Sigma; \Gamma \vDash A.$

Lemma

$\Gamma \models_g A \Rightarrow \Gamma; \Gamma \models A.$

Lemma (Algebraic compactness)

If $\Sigma; \Gamma \models A$, then there is a finite subset Γ_0 of Γ such that $\Sigma; \Gamma_0 \models A$.

Corollary

A Magari algebra is \square -founded if and only if it is embeddable into a σ -complete Magari algebra.

Proof sketch

Any \square -founded Magari algebra can be embedded into a reduced product of its countable (and finite) subalgebras over a countably complete filter.

Any countable (and finite) \square -founded Magari algebra is can be embedded into a σ -complete Magari algebra.

Any reduced product of σ -complete Magari algebras over a countably complete filter is σ -complete.

It seems that, relying only on DC without LEM, global completeness of KM over δ -complete KM-algebras can be obtained.

Proposition

A KM-algebra is \square -founded if and only if it is embeddable into a σ - δ -complete KM-algebra.