## Singly Generated Quasivarieties and Residuated Structures

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**Theorem** [essentially Maltsev 1966; Łoś & Suszko 1958]. The following conditions on a quasivariety **K** are equivalent.

(1)  $\mathbf{K} = \mathbb{Q}(\mathbf{A})$  for some  $\mathbf{A} \in \mathbf{K}$  (i.e.,  $\mathbf{K}$  is 'singly generated').

- (2) **K** has the *joint embedding property* (JEP), i.e., any two nontrivial members of **K** both embed into some third member.
- (3) For every set  $X \subseteq K$ , there exists  $B \in K$  s.t.  $X \subseteq \mathbb{IS}(B)$ .
- (4) [a robust 'relevance principle']:

For any finite set  $\Gamma \cup \Delta \cup \{\alpha \approx \beta\}$  of equations, where  $\Gamma$  is satisfiable in a nontrivial member of K and involves different variables from  $\Delta \cup \{\alpha \approx \beta\}$ ,

 $\text{if } \mathsf{K} \models (\&(\mathsf{\Gamma} \cup \Delta)) \Rightarrow \alpha \approx \beta, \text{ then } \mathsf{K} \models (\&\Delta) \Rightarrow \alpha \approx \beta.$ 

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As the JEP is a relevance principle, failing in **DMM**, our aim is to describe the (quasi)varieties of De Morgan monoids that have the JEP, or related 'completeness' properties.

(These algebras will be defined later.) It is profitable to conduct a universal-algebraic analysis first.

An algebra is said to be 0-*generated* if it has a distinguished element and no proper subalgebra.

Fact. In a quasivariety K with the JEP and a constant symbol,

- (1) the nontrivial 0-generated algebras are isomorphic, and
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**Theorem.** If **K** has the JEP, and if its nontrivial members lack trivial subalgebras, then there's a relatively simple  $A \in K$  such that  $\mathbb{ISP}_{\mathbb{U}}(A)$  includes every relatively simple algebra  $B \in K$ .

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# Although the JEP persists under category equivalences, it is not hereditary, i.e., it need not persist in sub(quasi)varieties.

Known variants of the JEP (with logically inspired names): **Definition.** A quasivariety **K** is

(i) structurally complete (SC) if  $\mathbf{K} = \mathbb{Q}(\mathbf{F}_{\mathbf{K}}(\aleph_0));$ 

(ii) *passively structurally complete* (PSC) if any two nontrivial members of K satisfy the same *existential positive sentences* ∃*x*<sub>1</sub> ... ∃*x*<sub>n</sub>Φ (Φ a disjunction of conjunctions of equations).
 Note. PSC is hereditary. Also, SC implies PSC.

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**Theorem.** Every PSC quasivariety has the JEP (hereditarily).

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The common retract C is finite and either relatively simple or trivial. A nontrivial common retract is unique, up to  $\cong$ .

E.g.: {groups}, {Heyting algebras}, {lattices}.

**Corollary.** If, moreover, the nontrivial members of **K** lack trivial subalgebras, then **K** has a unique relatively simple member.

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**Theorem.** The maximal PSC subquasivarieties of **DMM** are exactly the classes

 $\operatorname{Ret}(\operatorname{\mathsf{DMM}}, \operatorname{\boldsymbol{A}}) := \{ \operatorname{\boldsymbol{B}} \in \operatorname{\mathsf{DMM}} : \operatorname{\boldsymbol{A}} \text{ is a retract of } \operatorname{\boldsymbol{B}}, \text{ or } |\operatorname{\boldsymbol{B}}| = 1 \},\$ 

where **A** is any 0-generated De Morgan monoid.

Every nontrivial PSC quasivariety of De Morgan monoids is contained in just one of these.

Results of Slaney (1985) help to show that there are just 68 0-generated De Morgan monoids *A*, all of which are finite.

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It consists of idempotent subdirect products of chains.

It is generated by the chain  $\mathbb{Z}$  of all integers, in which e = 0,  $\neg x = -x$  and  $x \cdot y$  is whichever of x, y has the larger absolute value—or is  $x \wedge y$  if |x| = |y|.

Its subquasivarieties are varieties and they form a transparent chain of order type  $\omega + 1$ .

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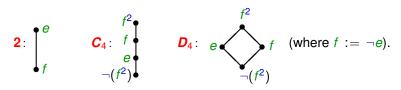
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For PSC/JEP in *varieties* of De Morgan monoids, we need: **Fact.** The simple 0–generated De Morgan monoids are just



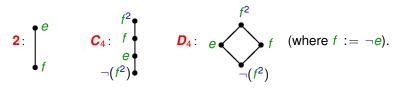
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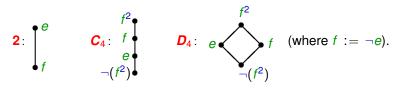


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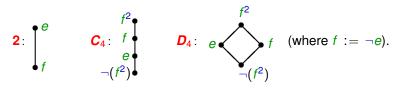
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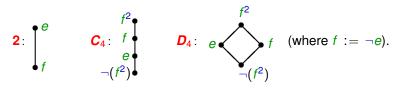
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**Claim.** M has  $2^{\aleph_0}$  (locally finite) subvarieties K that are structurally *incomplete*—i.e.,  $K = \mathbb{V}(L)$  for some proper subquasivariety L of K; in particular,  $\mathbb{Q}(F_K(\aleph_0)) \subsetneq K$ .

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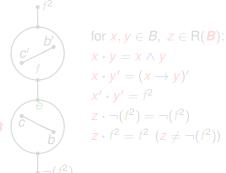
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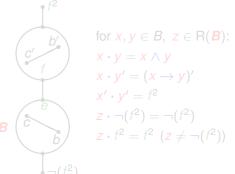
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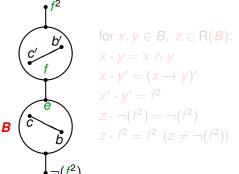
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Conversely, each Brouwerian algebra **B** has a '*reflection*'  $R(B) \in M$ , illustrated on the right [Meyer 1973]. for  $x, y \in B, z \in \mathbb{R}(B)$ :  $x \cdot y = x \wedge y$   $x \cdot y' = (x \rightarrow y)'$   $x' \cdot y' = f^2$   $z \cdot \neg (f^2) = \neg (f^2)$  $z \cdot f^2 = f^2 (z \neq \neg (f^2))$ 

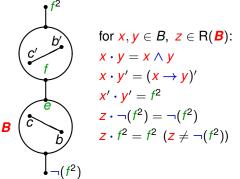
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Kuznetsov (1975) exhibited a set of  $\aleph_0$  finite Brouwerian algebras, distinct subsets of which generate distinct varieties.

His algebras can be modified so that, after a certain extra algebra **B** is added to the set, distinct subsets *containing* **B** generate distinct *structurally incomplete* subvarieties (that are locally finite).

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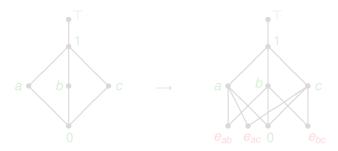
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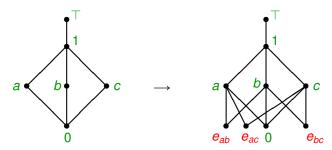
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For each doubleton  $\{x, y\}$  at depth 3, a new element  $e_{xy}$  is added, which has no strict lower bound; its set of strict upper bounds is the upward closure of  $\{x, y\}$  in the original poset.

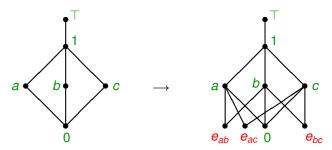
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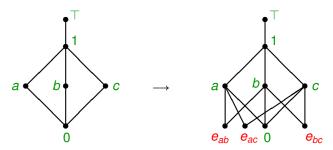
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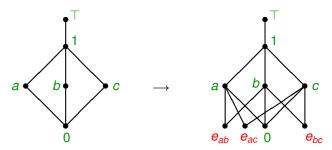
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